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(11)

EP 0 751 656 B1

(12)

EUROPEAN PATENT SPECIFICATION

(45) Date of publication and mention
of the grant of the patent:
16.08.2006 Bulletin 2006/33

(51) Int Cl.:
H04L 29/06 (2006.01)

(21) Application number: **96304706.3**

(22) Date of filing: **26.06.1996**

(54) **Protocol independent and adaptive network interface**

Protokollunabhängige, anpassungsfähige Netzwerkschnittstelle

Interface de réseau adaptatif et indépendant du protocole

(84) Designated Contracting States:
DE FR GB IT

(30) Priority: **27.06.1995 US 495172**

(43) Date of publication of application:
02.01.1997 Bulletin 1997/01

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Description

[0001] The present invention relates to a method of transmitting data between application programs executing on different devices interfaced to a local area network and, more particularly, to a method of transmitting data between application programs independent of any specific protocol.

[0002] Computerized local area networks (LAN's) are in widespread use for interconnecting many different computers and peripherals so as to allow users of the computers to communicate with one another and also to allow those users shared access to the peripherals. Recent developments in LAN's have seen the introduction of so-called "heterogeneous" LAN's, i.e., LAN's on which many different communication protocols are carried on a single Ethernet or Token-ring medium. Examples of different protocols are IPX, which is typically used by DOS-based PC's, UDP/IP, which is typically used by UNIX-based workstations, and DDP, which is typically used by Macintosh computers. Each type of computer or workstation can be adapted through software to communicate using multiple different protocols.

[0003] A peripheral also can include software, i.e., multiple protocol stack modules, which allows the peripheral to communicate using multiple protocols in order to be shared on a heterogeneous LAN. A protocol stack is a software module that processes packets of data which are received from or are transmitted to the LAN using the corresponding protocol. The protocol stacks and the associated lower-level software for network communications are typically stored and executed on a network interface device which may be embedded in or attached to the peripheral. The network interface device serves as an interface which allows the peripheral to communicate with other network devices via the LAN.

[0004] Network devices, i.e., network interface devices and computers which are interfaced to the LAN, also execute application programs. These programs execute at a level above the protocol stacks and can include, for example, print server programs, management programs which allow communication between a computer and a peripheral in order to configure or obtain status data from the peripheral, and other programs which may communicate data between different devices interfaced to the LAN. Exemplary management programs include programs that implement SNMP (Simple Network Management Protocol) and CMIP (Common Management Information Protocol). More than one such management program may be executing on a single network device at the same time.

[0005] US-A-5301303 discloses a network component in the form of a multiple channel concentrator which provides a number of physical communication channels for connection to LANs having different physical media such as twisted pair, coaxial or fibre optic and/or communication protocols such as Ethernet, fibre distribution data interface (FDDI), token bus and token ring. A number of

media distribution modules effective to implement a particular protocol and having ports for connection to a specific medium provide a connection between the particular medium and the connector. Thus, a network may be linked to the concentrator via a media distribution module to establish connection to another network. Bridging and routing modules within the concentrator allows communication between different channels and between networks having different protocols, either converting one protocol to another or filtering transmissions.

[0006] EP-A-0288713 discloses an input/output controller coupling a host computer to a device in an information processing system. Each input/output device has an associated control table for relating the device to program instructions stored in the controller defining communications protocol. When data is to be transferred between the input/output device and the host computer, information in the control table is used to call program instructions to impose the correct communications protocol to govern the transfer of the data to the input/output device.

[0007] In a conventional approach, an application program communicates with a protocol stack via an application programming interface (API) and uses the protocol stack to perform communications services. An application program must use different APIs to interface with each different protocol stack. This means that the application program must be aware of the particular network environment, i.e., the protocol in use, and the specific network API to be used.

[0008] The conventional approach leads to many difficulties. If an application program must support multiple network protocols, duplicated effort is required for the application software to handle the different APIs. For example, an SNMP program must include software code to communicate with an IPX protocol stack and a DDP protocol stack, in addition to code to communicate with a UDP protocol stack. Moreover, a CMIP program or any other application will also require the extra code for communicating with different protocol stacks. As a result, application programs that are designed to support multiple protocols using the conventional approach require a more complex design, have a longer development time, have less portability, and have a higher maintenance cost than application programs which support a single protocol.

[0009] Accordingly, a way is needed for application programs to communicate with application programs on other network devices independent of a specific protocol.

[0010] The above need is addressed by the present invention in which data is transmitted between application programs executing on different devices independent of a specific protocol.

[0011] According to the present invention, there is provided a method of delivering a data packet received from a first application program executing on a first device which is interfaced to a Local Area Network to a second application program executing on a second device which

is interfaced to the Local Area Network, the method comprising the steps of:

receiving a protocol independent data packet from the first - application program, together with data identifying the first application program and a destination ID identifying the second application program; determining which protocols are available for use on the Local Area Network; and assigning a respective access line to each of the protocols available for use on the Local Area Network, said method being characterised by the steps of:

selecting, from the available protocols, the protocol assigned to the access line having the least traffic;
determining protocol specific information which includes a type of protocol header and address information in the header corresponding to data identifying the first application program and the protocol selected in said selecting step;
forming a transmission packet including the data packet, the destination ID, and the determined protocol specific information; and
transmitting the transmission packet to the second application program via the Local Area Network.

[0012] By virtue of this arrangement, a data packet which is received from one application program without any data specifying a protocol can be transmitted to another application program based only on information identifying the source and destination programs received from the one application program.

BRIEF DESCRIPTION OF THE DRAWINGS

[0013]

Figure 1 is a diagram of a local area network.
Figure 2 is a diagram showing the software architecture used for communication between application programs according to an exemplary embodiment of the present invention.
Figure 3 is a diagram showing the functional relationships between software modules executing on a computer and a network interface device.
Figure 4 is a functional block diagram of a network expansion board for interfacing a printer to a local area network.
Figure 5 illustrates software modules that may be stored in memory on the network expansion board.
Figure 6 is a flow diagram showing process steps for transmitting a data packet between a first application program executing on a first network device and a second application program executing on a second network device.

Figures 7A and 7B show examples of protocol mapping tables created by protocol independent interfaces.

Figures 8A and 8B show examples of access ID mapping tables created by protocol independent interfaces.

Figures 9A and 9B show examples of protocol address mapping tables created by protocol independent interfaces.

Figures 10A through 10C show examples of transmission packet formats for use with different protocols.

Figure 11 is a flow diagram showing process steps for receiving a data packet transmitted from a first application program executing on one network device to a second application program executing on another network device.

DESCRIPTION OF AN EXEMPLARY EMBODIMENT

[1. System Overview]

[0014] Figure 1 is an illustration of a heterogeneous network system including several different types of computers and several different peripherals to which the computers can share access. The present invention can also be used with devices connected to a homogenous network, i.e., a network in which every device uses the same protocol.

[0015] In Figure 1, LAN 10 is depicted as an Ethernet medium which has a bus-type architecture, but a Token-ring medium having a ring-type architecture can be used as well. Connected to LAN 10 are a PC 20 which serves as a system administrator's computer, a PC 30 which serves as a print server for printers 85 and 95, a Macintosh computer 40, a UNIX workstation 50, and a generalized workstation 60 having a control unit 61 and a display 62. A fileserver 70 allows shared access to a network disk 75. A network expansion board (NEB) 100 allows shared access to a printer 105, and a network expansion device (NED) 110 allows shared access to a printer 115. In addition, a network interface board (NIB) 120 allows shared access to a copier 135 via a multiple device controller (MDC) 130.

[0016] The present invention relates to communication between application programs executing on different network devices. A preferred form of the present invention is described below in the context of communication between PC 20 and NEB 100. However, the present invention is applicable to computers and embedded network devices in general. Accordingly, the present invention can be applied to communication between other computers such as Macintosh computer 40, UNIX workstation 50, generalized workstation 60 and other network interface devices such as NED 110 (an example of which is described in copending U.S. Patent Application S.N. 08/489,116 filed on June 9, 1995, and entitled "Outputting a Network Device Log File") and NIB 120 (an example

of which is described in U.S. Patent Application S.N. 08/409,034 (corresponding to European application no. 96301994.8), filed on March 23, 1995, and entitled "Network Interface Board For Digital Copier", which is assigned to the assignee of the present invention).

[0017] The present invention also can be applied to communication between application programs executing on different computers that have the capacity to use multiple protocols. Moreover, the present invention is not limited to network applications, but instead can be used for devices having a direct connection through any bidirectional interface, e.g., a shared memory, a SCSI interface, an RS-1284 parallel interface, or the like.

[2. Software Architecture]

[0018] Figure 2 shows the software architecture of program modules executing on a network device such as NEB 100. Similar software executes on a computer such as PC 20. A network interface driver 210 is the lowest level of software which interfaces with LAN 10 and handles sending and receiving of packets on LAN 10 by adding or stripping off packet frame headers. Network interface driver 210 includes a media-type specific component which is designed for either an Ethernet or a Token-ring network medium and includes a plurality of logical boards for respectively processing packets having different frame types. A multiplexer software module 220 serves as a multiplexer which routes packets between network interface driver 210 and one or more protocol stacks. Each protocol stack receives packets that use the corresponding protocol, determines what needs to be done with the packets, and routes the packets to the appropriate application programs for servicing. The preferred embodiment supports three protocol stacks: IPX stack 230, UDP stack 231, and DDP stack 232. All of the protocol stacks might not be loaded on a particular device. Further, additional stacks for other protocols may be included. In the preferred embodiment, network interface driver 210, multiplexer software module 220, and protocol stacks 230 - 232 conform to the Open Data-Link Interface (ODI) specification described in "Open Data-Link Interface Developer's Guide for DOS Workstation Protocol Stacks", Version 1.10, Released by Novell, Inc., March 18, 1992.

[0019] A protocol independent interface (PII) 250 serves as an interface between protocol stacks 230 - 232 and management application programs such as an SNMP application program 260 and a CMIP application program 270. PII 250 "listens" for data packets addressed to particular sockets, i.e., addresses, and accepts those packets from the protocol stacks 230 - 232 for processing and forwarding to the application programs. Since SNMP program 260 and CMIP program 270 are executing on an embedded device, i.e., NEB 100, those programs are "agent" programs. An agent program collects and stores data regarding the network interface device, i.e., NEB 100, and the peripheral, i.e.,

printer 105, and responds to commands sent using the associated network management protocol, e.g., SNMP or CMIP, from a related "manager" program executing on a computer.

[0020] For example, the following predetermined addresses are used in the preferred embodiment for receiving data packets using the SNMP network management protocol: (1) for IPX, "socket" 900F_H and 9010_H (agent socket and trap socket, respectively), (2) for UDP, "port" 160_H and 161_H, and (3) for DDP, a unique name "SNMP Agent" and "SNMP Trap Handler" and "socket" 8_H and 9_H (agent socket and trap socket, respectively). Data packets that are not addressed to a socket used by a management program are routed to other application programs. For example, a PSERVER program 245 receives data packets via an API 240. Other application programs, including other management programs, can also be included.

[0021] Figure 3 is a diagram showing the functional relationship between manager programs executing on PC 20 and agent programs executing on NEB 100. As shown in Figure 3, PC 20 includes a protocol independent interface (PII) 255, an SNMP manager 265, and a CMIP manager 275. During an initialization process described below, PII 255 assigns a unique identifier referred to as an access ID to each management program. The access ID may be, for example, the MAC address of the device together with an additional number to uniquely identify each of SNMP manager 265 and CMIP manager 265. Reference numerals 281, 282, and 283 designate logical channels used by different respective protocols, such as IPX, UDP, and/or DDP. PII 255 assigns an identifier referred to as a "logical access line" to each of the protocols, i.e., logical channels 281 283, during initialization. By dynamically assigning access IDs and logical access lines, the PII can be readily adapted to support later-developed protocols without affecting the existing functionality.

[0022] PII 250 in NEB 100 likewise assigns an access ID to each of SNMP agent 260 and CMIP agent 270 and assigns a logical access line to each of the protocols. However, since the access IDs and logical access lines are specific to the PII that assigns them, the access IDs and logical access lines assigned by PII 250 in NEB 100 can differ from those assigned by PII 255 in PC 20. For example, as shown in Figure 3, SNMP manager 265 has access ID #2 in PC 20, but the corresponding agent in NEB 100, i.e., SNMP agent 260, has access ID #1. Similarly, as shown in Figure 3, PII 250 assigns logical access line #1 to logical channel 281, which may correspond to an IPX protocol, for example, while PII 255 assigns logical access line #3 to the same logical channel/protocol.

[0023] PII 250 and PII 255 perform identical interface functions. However, there are some differences in implementation due to the different platforms on which these software modules execute. For example, since PII 250 executes on an embedded device, it is implemented as a TSR (terminate and stay resident) routine. On the other hand, SNMP manager 265 is a Windows based applica-

tion and PII 255 is implemented as a Windows-based DLL (Dynamic Link Library). Other differences are indicated in the description below where appropriate.

[3. NEB Architecture]

[0024] Figure 4 is a functional block diagram of NEB 100. Broadly speaking, NEB 100 is an interactive network circuit board which couples printer 105 to LAN 10, making printer 105 a responsive and interactive network member. NEB 100 includes a shared memory SRAM 200 which is used for bidirectional communications between NEB 100 and printer 105. Printer 105 includes a printer interface card 220 (not shown) having a microprocessor 225 (not shown) that reads data from and writes data to SRAM 200. Printer 105 also includes a printer engine 250 (not shown) connected to printer interface card 220.

[0025] NEB 100 receives print data, status requests, and control commands from LAN 10, transmits print data, status requests, and control commands to printer 105 for execution, and transmits status information back to LAN 10. Thus, NEB 100 can perform not only RPRINTER remote printer services and PSERVER print server functionalities, but can also offer to network members whatever status and control features are available from the peripheral interface.

[0026] Power for all circuits is supplied to NEB 100 from a +5V power source 398. Power is provided from power source 398 to power converter 396 which provides -9V power to a transceiver 390 and to power converter 397 which provides +12V power to a flash EPROM 350 for "flashing" (i.e., reprogramming of the EPROM). Network and network interface control logic 340 is preferably a single 144-pin application specific integrated circuit (ASIC) that includes a network controller 330 and interface control logic 320. Network controller 330 is an NCR macro-cell compatible with a National DP83902A "ST-NIC" Ethernet controller, the details of which can be found in National Semiconductor's Local Area Networks Databook, National Semiconductor p/n 400055, National Semiconductor, 1993. Network controller 330 is designed to interface with CSMA/CA-type (carrier sense multiple access with collision detection) local area networks.

[0027] Network controller 330 connects with RJ-45 connector 385 directly and with coaxial connector 395 through transceiver 390, which is preferably a National Semiconductor DP8392 coaxial transceiver interface, the details of which can also be found in National's Local Area Networks Databook. Network controller 330 is also coupled to an 8KB SRAM 380 that is used as an input/output packet buffer for Ethernet data. This memory should preferably have an access time of about 70 ns or less.

[0028] Interface control logic 320 provides an interface between network controller 330, microprocessor 300, and memory devices EPROM 350 and DRAM 360. Interface control logic 320 also interfaces with non-volatile

random access memory (NVRAM) 370, which is a 256 byte serial electrically erasable/programmable memory used for initialization data storage during power cycling of printer 105. Network and printer configuration parameters are written into NVRAM 370 when printer 105 is first installed onto the network to allow NEB software to recover the installation parameters after printer power has been cycled off and on.

[0029] Interface control logic 320 also couples with serial port connector 325, which comprises a receive data pin 326 and a transmit data pin 327 that can respectively receive and transmit serial data streams for debugging purposes. Interface control logic 320 senses data present at the receive data line and samples the serial bits at regular intervals.

[0030] The central controller of NEB 100 is microprocessor 300, which is preferably an Intel 80C188EA-20 8-bit processor, the details of which can be found in the 80C186EA/80188EA User's Manual, Intel p/n 270950-001, Intel Corp. This processor is an 8-bit processor with direct memory access (DMA), interrupts, timers, and a DRAM refresh control. Other microprocessors, such as an AMD 80C188-20 8-bit microprocessor, might alternatively be used. 256 KB flash EPROM 350 and 512 KB DRAM 360 are coupled to microprocessor 300 via interface control logic 320, while 32 KB SRAM 200 (which is shared with printer interface card 220) is coupled with microprocessor 300 via arbiter control logic 400. A 40 MHz, 50 ppm crystal oscillator 310 provides microprocessor 300 with a clock signal that is wholly separate from and asynchronous with the clock signal provided to microprocessor 225 on printer interface card 220.

[0031] Microprocessor 300 executes instructions in flash EPROM 350, which stores control firmware and printing application software. After power-on self-test (POST), code from EPROM 350 is selectively moved to the higher performance 512 KB DRAM 360, which should preferably have an access time of about 80 ns, for actual execution.

[0032] All communication between NEB 100 and printer interface card 220 is executed via 32 KB shared SRAM 200. Arbiter control logic 400, preferably a single 100-pin ASIC, arbitrates between the two-byte-wide memory accesses of printer interface microprocessor 225 and the single-byte-wide memory accesses of NEB microprocessor 300, each of which is completely independent of the other.

[0033] Generally speaking, the 8-bit data bus of microprocessor 300 on board NEB 100 communicates with bus control logic 410, while the 32-bit data bus of microprocessor 225 on board printer interface card 220 communicates with bus control logic 420. Memory accesses from each bus are routed to shared memory arbiter 430, which determines which bus has priority and permits the bus with priority to access SRAM 200 via SRAM interface 440. Interrupt control register 450 is also accessed through shared memory arbiter 430, to allow one microprocessor to interrupt the other.

[0034] All software modules executed by microprocessor 300 are stored in flash EPROM 350. Those modules that are needed are selectively loaded from EPROM 350 into DRAM 360 and are executed from DRAM. This permits flexible configuration of NEB 100 by selection of which modules are to be loaded.

[0035] Figure 5 illustrates an example of blocks of code, or software modules, that are stored in flash EPROM 350. The PII module contains process steps for providing the required functions of a protocol independent interface, as described in more detail below. The XPL module provides a standardized interface between printer 105 and NEB 100. MLID (Multi Link Interface Driver) serves as network interface driver 210 and is a piece of Novell code (Media Support Module, or MSM) linked together with a piece of customized code (Hardware Support Module, or HSM) that is the lowest level of network connection, while LSL (Link Support-Layer) serves as multiplexer software module 220 and is a piece of Novell code that acts as a multiplexer between the low level MLID and the several protocol stacks above it. CNETX is customized code that turns local DOS-like function calls into network function calls, providing file functions like OPEN, READ, WRITE, and CLOSE.

[0036] The PRETASK module is responsible for identifying what frame types are associated with the various possible protocol stacks. Because NEB 100 supports multiple protocol stacks, this module exists as long as NEB 100 is running.

[0037] Novell's IPX/SPX protocol stack is contained in flash EPROM 350, and is supported by SAP, or Service Advertising Protocol. SAP is a Novell concept that allows devices to register themselves into the file server's bindery, which lists active and inactive network entities. Because print servers are a special kind of bindery item, SAP registers NEB 100 via CPSOCKET, and if NEB 100 is configured as a print server, SAP also registers the print server with the NetWare bindery.

[0038] CPSEVER is a custom implementation of a Novell print server application. This module provides self-generated print banners, user notification of completion and exception status, and transmission of print data and status commands to the printer. This differs from the Novell print server in that CPSEVER is dedicated to driving the local printer (i.e., printer 105 in which NEB 100 is installed) and cannot drive any remote RPRINTERS. This program owns the print data lines for the duration of a print job. CRPRINTER is a custom implementation of a Novell RPRINTER print application. This module is a slave application that is sent data by a Novell print server application elsewhere on LAN 10.

[0039] The TCP/IP protocol stack has User Datagram Protocol (UDP), Reverse Address Resolution Protocol (RARP) and BootP support within. INTERRUPT is the interrupt handler for the TCP/IP task, while TIMERTICK is the timer tick for UNIX TCP/IP network tasks. LPRINTER is the TCP/IP print server application, and also owns the print data lines for the duration of a print job.

DDP is the module for implementing a Datagram Delivery Protocol (DDP) which is used, for example, for communications with a Macintosh computer.

[0040] The CPSOCKET program runs for all protocol stacks. The program responds to requests for Connection, requests for data download, or requests for services from remote utilities, and provides status and control to other tasks via interprocess communication. Because CPSOCKET typically owns the status and control lines between NEB 100 and printer 105, it is the only task that has the ability to obtain printer status via the status lines. CPSOCKET is responsible for the network connection and packet contents between the Novell-oriented status and control utilities (CPNET or the corresponding Windows version of client-based software utilities), or between the UNIX-oriented status and control utilities (CPUTIL).

[0041] MONITOR is a customized multi-tasking monitor which performs task creation, task destruction and microprocessor dispatch. MONITOR also has memory management sub-modules MEMGET and MEMFREE. RESIDENT is a block of routines that provides generic services such as read and write to flash EPROM 350, FLASH code, ROM based debugger, hardware timer tick and other basic features. POST is a power-on self-test module that checks the integrity of NEB hardware and software at power-up.

[0042] Also stored in EPROM 350 is a network identification file (NIF) data block which stores board-invariant information, which is unique for every network board, hardware configuration data, board revision number and the like, as well as changeable information such as software version number. The information in the NIF data block is used to ensure that flash EPROM 350 is not reprogrammed with an incompatible firmware image.

[0043] Specifically, EPROM 350 stores "board" information such as model number, firmware level, and board revision number, as well as "network" information such as Media Access Control (MAC) address, which is unique for every network board, board name, network frame type, primary file server identification, queues serviced, network protocol, sampling frequency, PSEVER name, zone-name, and the like.

[4. Function of PII]

[0044] Figure 6 is a flow diagram showing process steps for implementing a protocol independent method of transmitting a data packet from a first application program executing on a first device which is interfaced to a LAN to a second application program executing on a second device which is interfaced to the LAN. Briefly, according to Figure 6, a protocol independent interface (PII) program is initialized which determines which protocols are available for use, assigns an access line to each protocol that is available for use, assigns an access ID to the first application program, and creates mapping information that indicates a one-to-one correspondence be-

tween an access ID/access line pair and a block of protocol specific information which includes a protocol header having predetermined address data. This one-to-one mapping can be done by using a mapping table, as in the preferred form described below, or by implementing a data structure that carries the mapping information. A data packet is sent to the PII program together with the access ID of the first application program and a destination ID for the second application program, and one of the available protocols is selected to transmit the data packet. A block of protocol specific information is retrieved from the mapping table based on the access ID of the first application program and the access line corresponding to the selected protocol, and a transmission packet is formed which includes the data packet, the destination ID, and the retrieved block of protocol specific information. The transmission packet is then transmitted via the LAN.

[0045] More specifically, the process steps of Figure 6 show transmission of a data packet from SNMP manager 265 on PC 20 to SNMP agent 260 on NEB 100. The function of PII 250 to transmit a data packet from SNMP agent 260 to SNMP manager 265 is the same, with the only differences being the differences in implementation discussed herein. In step S601, SNMP manager 265 executes an initialization command to initialize PII 255. This command must be executed prior to execution of any other PII commands and need only be performed once. As mentioned above, PII 255 is implemented in Windows as a DLL. Accordingly, in response to the initialization command, the necessary operations are performed to enable SNMP manager 265 to execute other PII commands using the DLL. In contrast, initialization of PII 250 by SNMP agent 260 in NEB 100 returns a table of entry points that SNMP agent 260 uses to call other PII routines.

[0046] PII 255 then determines which protocols are available for use by PC 20. In the preferred embodiment, where PC 20 is operating in a Windows or DOS Novell-ODI environment, function calls which obtain an indication of the presence or absence of each protocol stack are issued in a round-robin manner to determine which protocol stacks are available. Alternatively, commands to execute respective initialization routines for each protocol stack can be issued in round-robin fashion. If an initialization routine fails, the failure is interpreted to indicate that the protocol stack is not available at that instant. In the embedded platform, i.e., for PII 250 on NEB 100, the pretask module has information regarding the available protocol stacks and that information is obtained and used by PII 250.

[0047] After determining which protocols are available, PII 255 opens a socket for each available protocol depending on the type of protocol, and sets up a protocol mapping table. This table has a one-to-one correspondence between an access line and a type of protocol stack. As noted above, the mapping table is one of many possible implementations. For example, the protocol map-

ping also can be implemented as a bit map or some other data structure. The key is that there is a way of indicating a one-to-one correspondence between the access line and protocol stack. The mapping table used in the preferred embodiment is formed by assigning an access line to each available protocol and storing data indicating the access line assigned to each protocol in a section of memory reserved for use by PII 255. Figure 7A is an example of a protocol mapping table that may be created by PII 255 when UDP, IPX, and DDP protocols are available, using the exemplary access line assignments shown in Figure 3.

[0048] After initialization of PII 255 in step S601, flow proceeds to step S602 in which an Open command is executed by SNMP manager 265 to open a session. This command returns an access ID for a management program to use to identify itself, e.g., access ID #2 for SNMP manager 265. PII 255 stores data indicating the relationship between access IDs and management programs in the section of memory reserved for use by PII 255. Figure 8A is an example of an access ID mapping table set up by PII 255 to indicate the relationship between access IDs and management programs. As with the protocol mapping table discussed above, the mapping between access IDs and application programs is capable of many implementations other than a mapping table, e.g., a bit map or other data structure. The data in Figure 8A corresponds to the exemplary assignment of access IDs shown in Figure 3. The access IDs are shown as XXXX1 or XXXX2, where XXXX represents the MAC address of PC 20 and the 1 and 2 indicate CMIP manager 275 and SNMP manager 265, respectively. The purpose of an access ID is to uniquely identify different entities that utilize the PII, for example, CMIP manager 275 and SNMP manager 265.

[0049] Flow next advances to step S603 in which a data packet is sent to PII 255 from SNMP manager 265. The packet is sent by providing a pointer to the packet's location in memory, together with the access ID for SNMP manager 265 and a destination ID for the packet's destination. Since no information indicating a transmission protocol is required for the packet, the packet is protocol independent and a single interface is provided between an application program and all protocol stacks. The destination ID of the corresponding agent, e.g., SNMP agent 260 corresponding to SNMP manager 265, is determined in the preferred embodiment by performing a locate_agent function before data is sent to an agent. For example, in the case where Novell Netware is used, this function is performed by looking in a bindery in fileserver 70 to determine devices which are registered in the bindery, e.g., by using IPX's SAP function, and which may have compatible agents. Communication with those devices then takes place to obtain a destination ID for the agent. In a UNIX environment, the locate_agent function is performed by using predetermined host tables. In an AppleTalk environment, an agent is located using a unique name. Any other technique can also be used

which allows the manager program to obtain a destination ID for an agent program.

[0050] After the packet is sent to PII 255, flow proceeds to step S604 in which a protocol is selected to transmit the data packet. If only one protocol is available, that protocol will be used. In the preferred embodiment, a protocol is selected by using a default or preferred protocol, if available. For example, UDP is the preferred protocol for transmitting data between SNMP manager 265 and SNMP agent 260. If the preferred protocol is not available, then the first available protocol is used. However, there are many alternative variations for selecting a protocol to use. For example, a protocol can be selected randomly or the first available protocol can be used. Also, the protocol having the least traffic can be used. For example, in a Novell environment, library functions such as GetLocalTarget can return an estimate of the time required to deliver a 576-byte packet to a designated destination. These functions provide information related to network traffic using the corresponding protocol. A call to those functions can be used to obtain information indicating which protocol has the least traffic. Further, PII 255 can store a counter for each protocol (or access line) and can select the one that PII 255 has used the fewest times, or PII 255 can store a time at which each protocol (or access line) was last used and can select the one that was used least recently.

[0051] After selecting a protocol to transmit the data packet, flow advances to step S605 in which PII 255 retrieves a block of protocol specific information, i.e., a protocol specific address, from a protocol address mapping table stored in memory. The information is retrieved based on the access ID of the transmitting program and the access line corresponding to the selected protocol. If the mapping table does not have an entry for the Access ID/Access line pair, for example, if the packet is the first packet transmitted by PII 255 for a particular manager using a particular protocol, an entry is added to the table.

[0052] Figure 9A illustrates an example of a protocol address mapping table created by PII 255. There is a one-to-one correspondence between an Access ID/Access line pair and a protocol specific address. Reference numeral 901 designates a column of Access ID/Access line pairs. Reference numeral 902 represents a column of protocol specific addresses that each include a protocol header of a type indicated in column 903. Each protocol header has a different format. The various protocol header formats are described in detail below with respect to Figures 10A - 10C. However, regardless of type, each protocol header includes some type of address data, e.g., a socket for an IPX header, a port for a UDP header, and a socket and name for a DDP header. Reference numeral 904 indicates a column showing the specific address data included in each protocol header.

[0053] The address data indicated in column 904 represents standard address data that is defined for use by a particular manager and protocol. For example, the SNMP protocol (access ID XXXX2 in Figures 8A and 9A)

uses the following address data: (1) for IPX, "socket" 900F_H and 9010_H (agent socket and trap socket, respectively), (2) for UDP, "port" 160_H and 161_H, and (3) for DDP, a unique name "SNMP Agent" and "SNMP Trap Handler" and "socket" 8_H and 9_H (agent socket and trap socket, respectively). Likewise, the CMIP protocol uses specific address data such as a CMIP port for UDP, a CMIP socket for IPX, and a CMIP name and socket for DDP. An SNMP protocol packet and a CMIP protocol packet may have the same header type, but each header will include different specific address data. For example, the protocol specific addresses shown in rows 1 and 4 of Figure 9A each have a UDP-type header, but each header contains different address data as shown in column 904. If SNMP manager 265 has access ID XXXX2, as indicated in Figure 8A, and the preferred protocol UDP is being used and has access line #1, as indicated in Figure 7A, then the protocol specific address in the fourth row of Figure 9A will be retrieved by PII 255, which corresponds to Access ID/Access line pair XXXX2/1.

[0054] If the present invention is used with application programs that do not have predefined sockets, standard identification values must be defined for those application programs and the PIs used with both manager and agent programs must be programmed to use those standard identification values.

[0055] As mentioned above, each type of protocol header has a different format, as described with respect to Figures 10A - 10C. Figure 10A shows a format for a local network frame (or "transmission packet") 1000 when the frame is transmitted using a UDP protocol. In this case, network frame 1000 includes a local network header 1010 and a local network trailer 1015. It also includes an IP header 1021, a UDP header 1020, and data 1030. As shown in Figure 10A, UDP header 1020 includes fields for a source port, a destination port, a length, and a checksum. Figure 10B shows a format for network frame 1000 when the frame is transmitted using an IPX protocol. In that case, network frame 1000 does not include IP header 1021 and includes an IPX header 1025 in lieu of UDP header 1020. IPX header 1025 includes fields for a checksum, a packet length, a transport control value, a packet type, a destination network, a destination node, a destination socket, a source network, a source node, and a source socket. Lastly, Figure 10C shows a format for network frame 1000 when the frame is transmitted using a DDP protocol. In this case, the network frame also lacks IP header 1021 and includes a DDP header 1028 instead of either UDP header 1020 or IPX header 1025. DDP header 1028 includes fields containing, 00, Hop, a datagram length, a datagram checksum, a destination network, a source network, a destination node ID, a source node ID, a destination socket number, a source socket number, and a DDP type.

[0056] Referring again to Figure 6, after retrieving the protocol specific information, flow proceeds to step S606 in which PII 255 forms a transmission packet which is specific to the selected protocol, i.e., a network frame

1000 having one of the formats shown in Figures 10A - 10C. The transmission frame is formed using the destination information provided by SNMP manager 265 and the information retrieved from the protocol address mapping table in step S605. Further, data 1030 is the data being sent by SNMP manager 265.

[0057] After forming transmission packet 1000, flow proceeds to step S607 in which transmission packet 1000 is transmitted to the destination application program via LAN 10.

[0058] Figure 11 is a flow diagram showing process steps for receiving transmission packet 1000 at NEB 100. In step S1101, SNMP agent 260 issues a command to initialize PII 250. PII 250 must be initialized before it can receive data from LAN 10 or SNMP agent 260. As mentioned above, SNMP agent 260 obtains a table of entry points into routines of PII 250 upon initialization, and PII 250 determines which protocols are available by obtaining information from the PRETASK software module. Figure 7B shows an exemplary protocol table that is set up and stored by PII 250 when UDP, IPX, and DDP protocols are available and access lines are assigned as shown in Figure 3.

[0059] Flow then advances to step S1102 in which SNMP agent 260 issues a PII open command to obtain an access ID. Figure 8B shows an exemplary mapping of access IDs to management agents that is set up and stored by PII 250, based on the exemplary assignment of access IDs shown in Figure 3. In Figure 8B, YYYY represents the MAC address of NEB 100. SNMP agent 260 is a passive entity which responds when SNMP manager 265 requests information but does not initiate communication. Accordingly, SNMP agent 260 listens for packets on all protocols that are available, i.e., packets addressed to any socket defined for SNMP use. Therefore, once an access ID is provided for an agent, e.g., SNMP agent 260, a mapping table is created with access ID/access line pair entries for every access line. Figure 9A shows an exemplary protocol address mapping table for PII 250 based on the exemplary data in Figures 7A and 8A.

[0060] Flow then advances to step S1103 in which a transmission packet is received from LAN 10. Referring to Figure 2, transmission packet 1000 is received from LAN 10 by network interface driver 210 and is routed to multiplexer software module 220 and then to one of protocol stacks 230 - 232. For example, when transmission packet 1000 represents data sent from SNMP manager 265 to SNMP agent 260 using the preferred UDP protocol, transmission packet 1000 is routed to UDP protocol stack 231. PII 250 listens for packets addressed to specific sockets, e.g., sockets defined for use by management programs, and ignores all others. Since transmission packet 1000 is addressed to one of the sockets to which PII 250 listens, PII 250 will receive the packet.

[0061] Flow then advances to step S1104 in which PII 250 obtains an Access ID/Access line pair based on protocol specific information 1020 in transmission packet

1000 by referring to the protocol address mapping table for PII 250 shown in Figure 9A. For exemplary transmission packet 1000, Access ID/Access line pair YYYY1/2 will be retrieved from row 2 in Figure 9A.

[0062] Flow then advances to step S1105 in which data 1030 is passed to the appropriate management program. This is done by referring to the access ID mapping information for PII 250, which is discussed above with reference to Figure 8B. For example, Access ID YYYY1 corresponds to SNMP agent 260, so data 1030 of transmission packet 1000 is passed to SNMP agent 260.

[0063] Although an example has been described for transmitting data from a manager in PC 20 to an agent in NEB 100, the same process is applicable for transmitting data from NEB 100 to PC 20. Moreover, as discussed above, the present invention is not limited to transmission of data between management application programs and is not limited to network transmissions. Accordingly, while the preferred embodiment of the invention has been described, it is to be understood that the invention is not limited to the above-described embodiments and that various changes and modifications may be made by those of ordinary skill in the art without departing from the scope of the invention.

Claims

1. A method of delivering a data packet received from a first application program executing on a first device which is interfaced to a Local Area Network (10) to a second application program executing on a second device which is interfaced to the Local Area Network (10), the method comprising the steps of:

receiving a protocol independent data packet from the first application program, together with data identifying the first application program and a destination ID identifying the second application program;
determining which protocols are available for use on the Local Area Network (10); and
assigning a respective access line (281, 282, 283) to each of the protocols available for use on the Local Area Network,
said method being **characterised by** the steps of:

selecting, from the available protocols, the protocol assigned to the access line having the least traffic;
determining protocol specific information which includes a type of protocol header and address information in the header corresponding to data identifying the first application program and the protocol selected in said selecting step;
forming a transmission packet including the

- data packet, the destination ID, and the determined protocol specific information; and transmitting the transmission packet to the second application program via the Local Area Network (10). 5
2. A method according to Claim 1, wherein the protocol independent interface program:
- assigns a unique access ID to the first application program, and 10
- creates mapping information having a one-to-one correspondence between each access ID/ access line pair (90) and a block of protocol specific information, and wherein said determining 15
- step comprises retrieving a block of protocol specific information from the mapping information based on the access ID of the first application program and the access line (281, 282, 283) 20
- corresponding to the protocol selected in said selecting step.
3. A method according to Claim 1 or 2, wherein the second device is a network interface device (7.20) 25
- which interfaces between a peripheral (135) and the Local Area Network (10) and which executes a service routine, wherein said method further comprises the step of transmitting data packets including data to be serviced by the service routine. 30

Patentansprüche

1. Verfahren zum Zuführen eines Datenpakets, das von einem ersten Anwendungsprogramm empfangen wird, das auf einer ersten Vorrichtung läuft, die eine Schnittstelle zu einem lokalen Bereichsnetzwerk (10) aufweist, zu einem zweiten Anwendungsprogramm, das auf einer zweiten Vorrichtung läuft, die eine Schnittstelle zu dem lokalen Bereichsnetzwerk (10) aufweist, wobei das Verfahren die Schritte umfasst:
- Empfangen eines protokollunabhängigen Datenpakets von dem ersten Anwendungsprogramm zusammen mit Daten, die das erste Anwendungsprogramm identifizieren, und einer Ziel-ID, die das zweite Anwendungsprogramm identifiziert, 45
- Bestimmen, welche Protokolle in dem lokalen Bereichsnetzwerk (10) zur Verwendung zur Verfügung stehen, und
- Zuweisen einer jeweiligen Zugriffsleitung (281, 282, 283) zu jedem der Protokolle, die in dem lokalen Bereichsnetzwerk (10) zur Verwendung zur Verfügung stehen, 55

wobei das Verfahren **gekennzeichnet ist durch** die

Schritte:

Auswählen, aus den zur Verfügung stehenden Protokollen, des Protokolls, das der den geringsten Verkehr aufweisenden Zugriffsleitung zugewiesen ist,

Bestimmen von protokollspezifischen Informationen, die eine Art eines Protokoll-Headers und Adressinformationen in dem Header enthalten, die Daten entsprechen, die das erste Anwendungsprogramm und das in dem Auswahlschritt ausgewählte Protokoll identifizieren,

Bilden eines Übertragungspakets, das das Datenpaket, die Ziel-ID und die bestimmten protokollspezifischen Informationen enthält, und Übertragen des Übertragungspakets über das lokale Bereichsnetzwerk (10) zu dem zweiten Anwendungsprogramm.

2. Verfahren gemäß Anspruch 1, wobei das protokollunabhängige Schnittstellenprogramm:

dem ersten Anwendungsprogramm eine eindeutige Zugriffs-ID zuweist, und

Abbildungsinformationen erstellt, die eine Eins-zu-Eins-Entsprechung zwischen jeder/jedem Zugriffs-ID/Zugriffsleitungspaar (90) und einem Block von protokollspezifischen Informationen aufweisen, und wobei der Bestimmungsschritt ein Ermitteln eines Blocks von protokollspezifischen Informationen aus den Abbildungsinformationen auf der Grundlage der Zugriffs-ID des ersten Anwendungsprogramms und der Zugriffsleitung (281, 282, 283), die dem in dem Auswahlschritt ausgewählten Protokoll entspricht, umfasst.

3. Verfahren gemäß Anspruch 1 oder 2, wobei die zweite Vorrichtung eine Netzwerkschnittstellenvorrichtung (120) ist, die eine Schnittstelle zwischen einer Peripherie (135) und dem lokalen Bereichsnetzwerk (10) bildet, und die eine Dienstroutine ausführt, wobei das Verfahren ferner den Schritt des Übertragens von Datenpaketen, die durch die Dienstroutine zu bedienende Daten enthalten, umfasst.

Revendications

1. Procédé pour délivrer un paquet de données reçues d'un premier programme d'application s'exécutant sur un premier dispositif qui est en interface avec un réseau local (10) à un second programme d'application s'exécutant sur un second dispositif qui est en interface avec le réseau local (10), le procédé comprenant les étapes qui consistent :

à recevoir un paquet de données indépendant

du protocole depuis le premier programme d'application, en même temps que des données identifiant le premier programme d'application et une ID de destination identifiant le second programme d'application ;
à déterminer les protocoles qui sont disponibles pour une utilisation sur le réseau local (10) ; et
à affecter une ligne d'accès respective (281, 282, 283) à chacun des protocoles disponibles pour une utilisation sur le réseau local,
ledit procédé étant **caractérisé par** les étapes qui consistent :

à sélectionner, parmi les protocoles disponibles, le protocole affecté à la ligne d'accès ayant le moindre trafic ;
à déterminer une information spécifique de protocole qui comprend un type d'en-tête de protocole et une information d'adresse dans l'en-tête correspondant à des données identifiant le premier programme d'application et le protocole sélectionné dans ladite étape de sélection ;
à former un paquet d'émission comprenant le paquet de données, l'ID de destination et l'information spécifique de protocole déterminée ; et
à transmettre le paquet d'émission au second programme d'application:en passant par le réseau local (10).

2. Procédé selon la revendication 1, dans lequel le programme d'interface indépendant du protocole :

affecte un ID d'accès unique au premier programme d'application, et
crée une information d'application ayant une correspondance biunivoque entre chaque ID d'accès/paire de lignes d'accès (90) et un bloc d'information spécifique de protocole, et dans lequel ladite étape de détermination comprend la récupération d'un bloc d'information spécifique de protocole à partir de l'information d'application sur la base de l'ID d'accès du premier programme d'application et de la ligne d'accès (281, 282, 283) correspondant au protocole sélectionné dans ladite étape de sélection.

3. Procédé selon la revendication 1 ou 2, dans lequel le second dispositif est un dispositif (120) d'interface de réseau qui établit une interface entre un périphérique (135) et le réseau local (10) et qui exécute un sous-programme de service, dans lequel ledit procédé comprend en outre l'étape d'émission de paquets de données comprenant des données devant être desservies par le sous-programme de service.

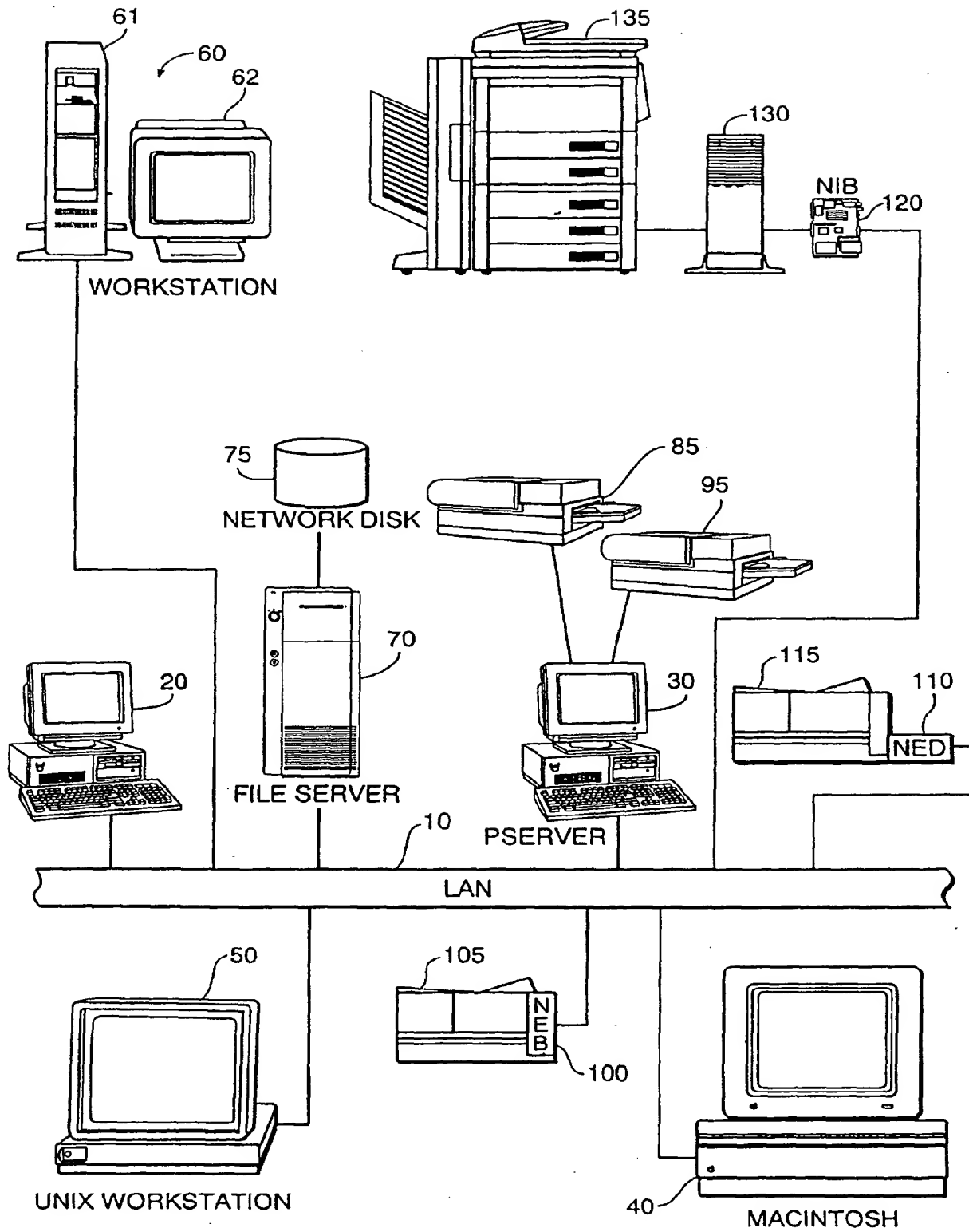
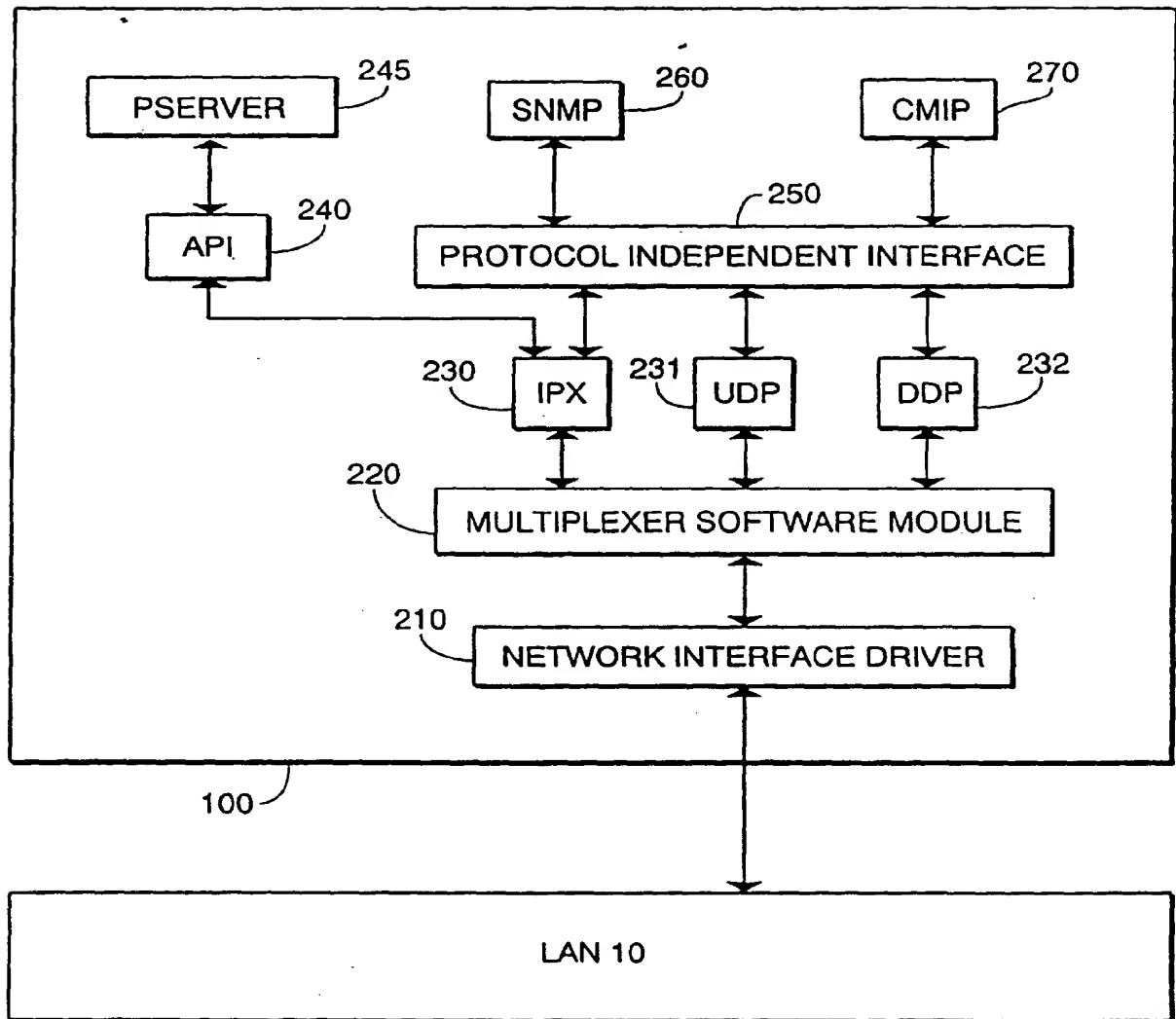


FIG. 1

**FIG. 2**

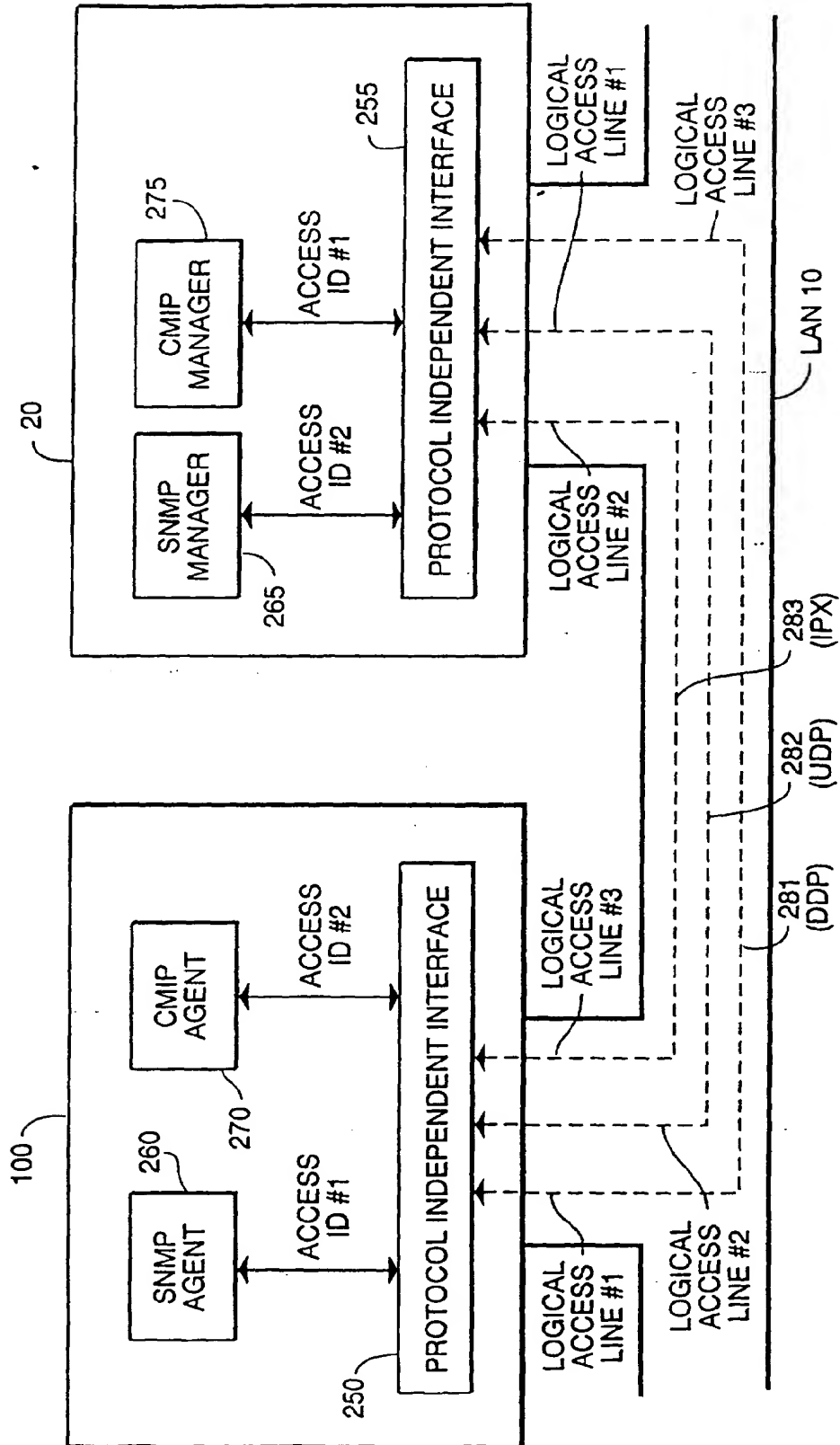


FIG. 3

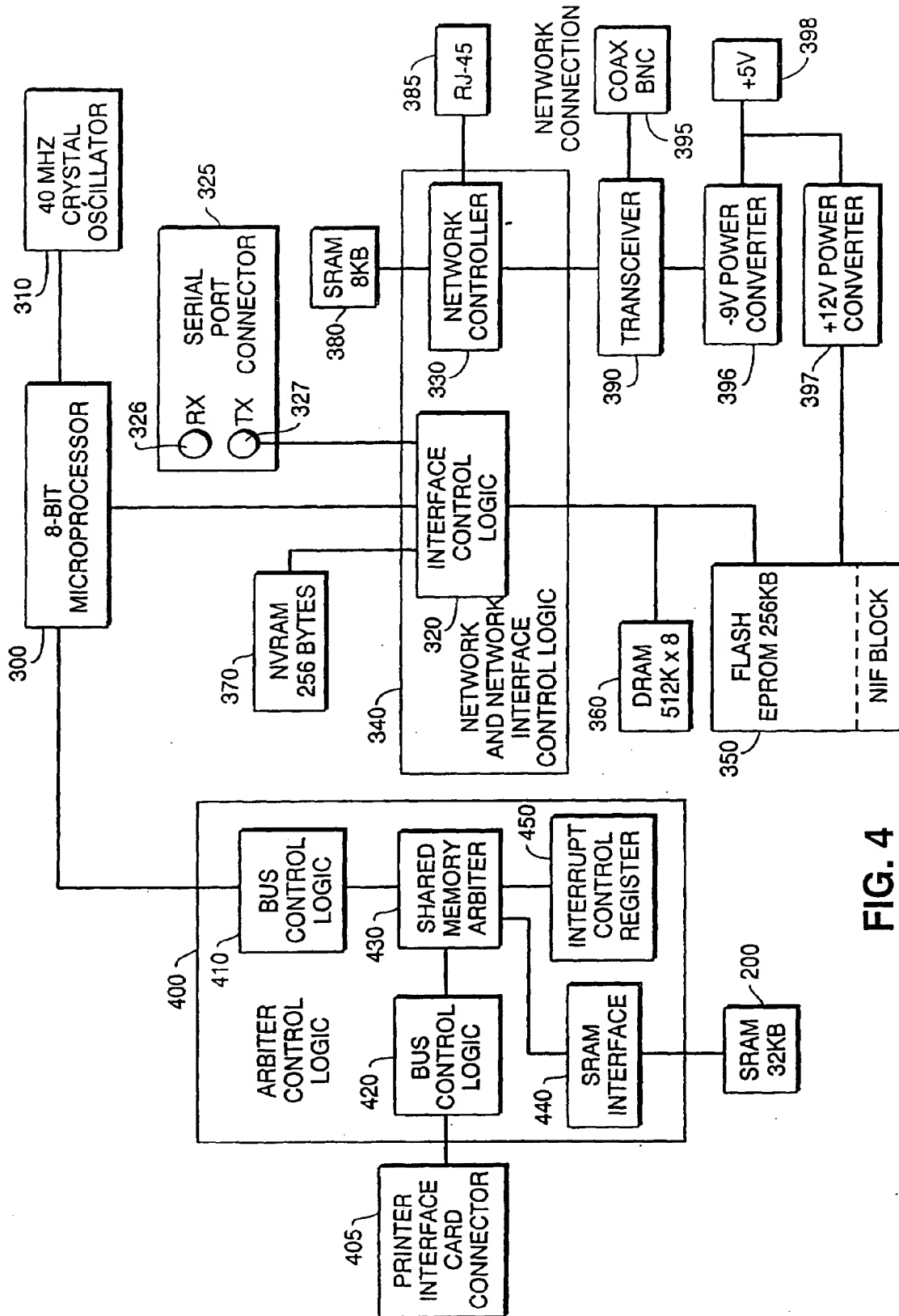
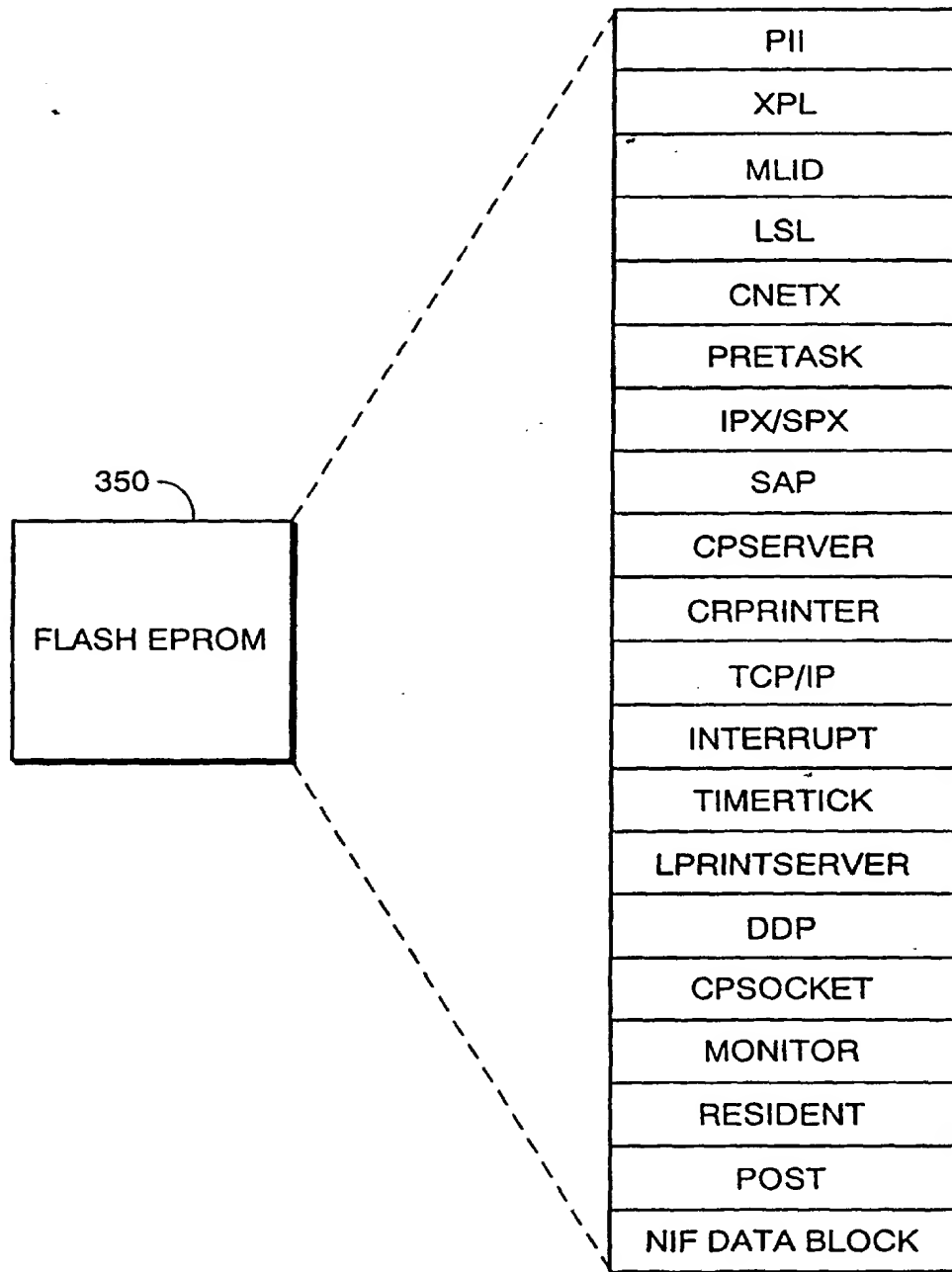


FIG. 4

**FIG. 5**

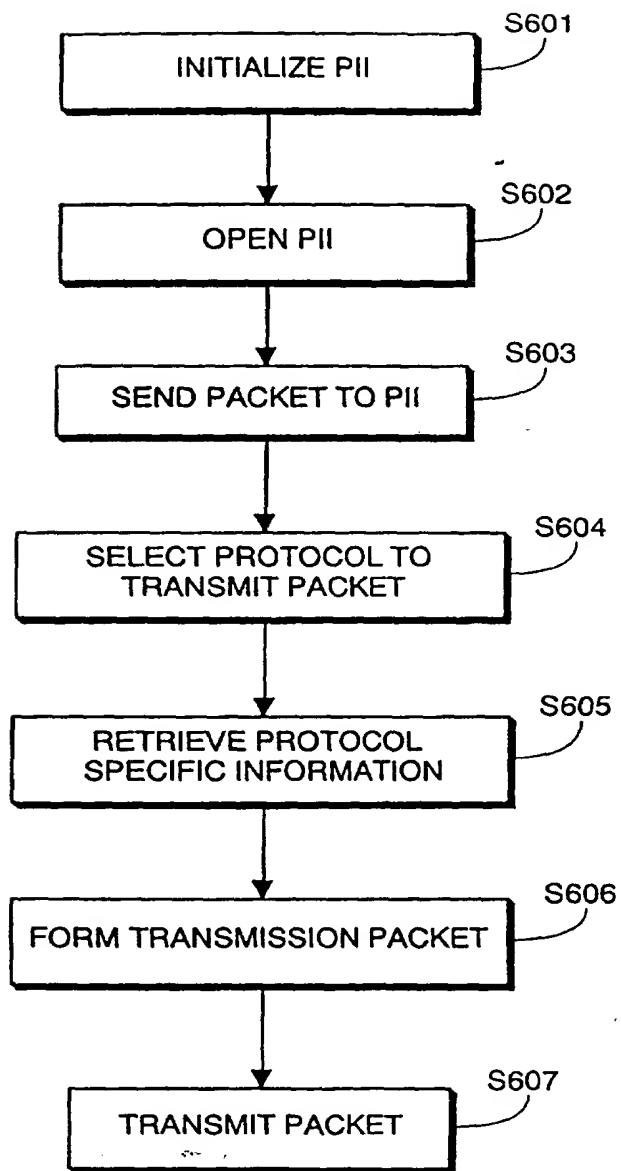


FIG. 6

FIG. 7A

<u>PROTOCOL</u>	<u>ACCESS LINE #</u>
UDP	1
IPX	2
DDP	3

FIG. 7B

<u>PROTOCOL</u>	<u>ACCESS LINE #</u>
UDP	2
IPX	3
DDP	1

FIG. 8A

<u>PROGRAM</u>	<u>ACCESS ID #</u>
CMIP MANAGER	XXXX1
SNMP MANAGER	XXXX2

FIG. 8B

<u>PROGRAM</u>	<u>ACCESS ID #</u>
CMIP AGENT	YYYY2
SNMP AGENT	YYYY1

FIG. 9A

<u>ACCESS ID/ACCESS LINE</u>	<u>PROTOCOL SPECIFIC ADDRESS</u>	
	<u>HEADER TYPE</u>	<u>ADDRESS IN HEADER</u>
XXXX1/1	UDP HEADER	CMIP PORT
XXXX1/2	IPX HEADER	CMIP SOCKET
XXXX1/3	DDP HEADER	CMIP SOCKET/NAME
XXXX2/1	UDP HEADER	PORT 160/161
XXXX2/2	IPX HEADER	SOCKET 900F/9010
XXXX2/3	DDP HEADER	SOCKET 8/9 - SNMP AGENT/SNMP TRAP HANDLER

901 → (points to XXXX2/3)

902 → (points to FIG. 9A)

903 → (points to DDP HEADER in row XXXX2/3)

904 → (points to SOCKET 8/9 - SNMP in row XXXX2/3)

FIG. 9B

<u>ACCESS ID/ACCESS LINE</u>	<u>PROTOCOL SPECIFIC ADDRESS</u>	
	<u>HEADER TYPE</u>	<u>ADDRESS IN HEADER</u>
YYYY1/1	DDP HEADER	SOCKET 8/9 - SNMP AGENT/SNMP TRAP HANDLER
YYYY1/2	UDP HEADER	PORT 160/161
YYYY1/3	IPX HEADER	SOCKET 900F/9010
YYYY2/1	DDP HEADER	CMIP SOCKET/NAME
YYYY2/2	UDP HEADER	CMIP PORT
YYYY2/3	IPX HEADER	CMIP SOCKET

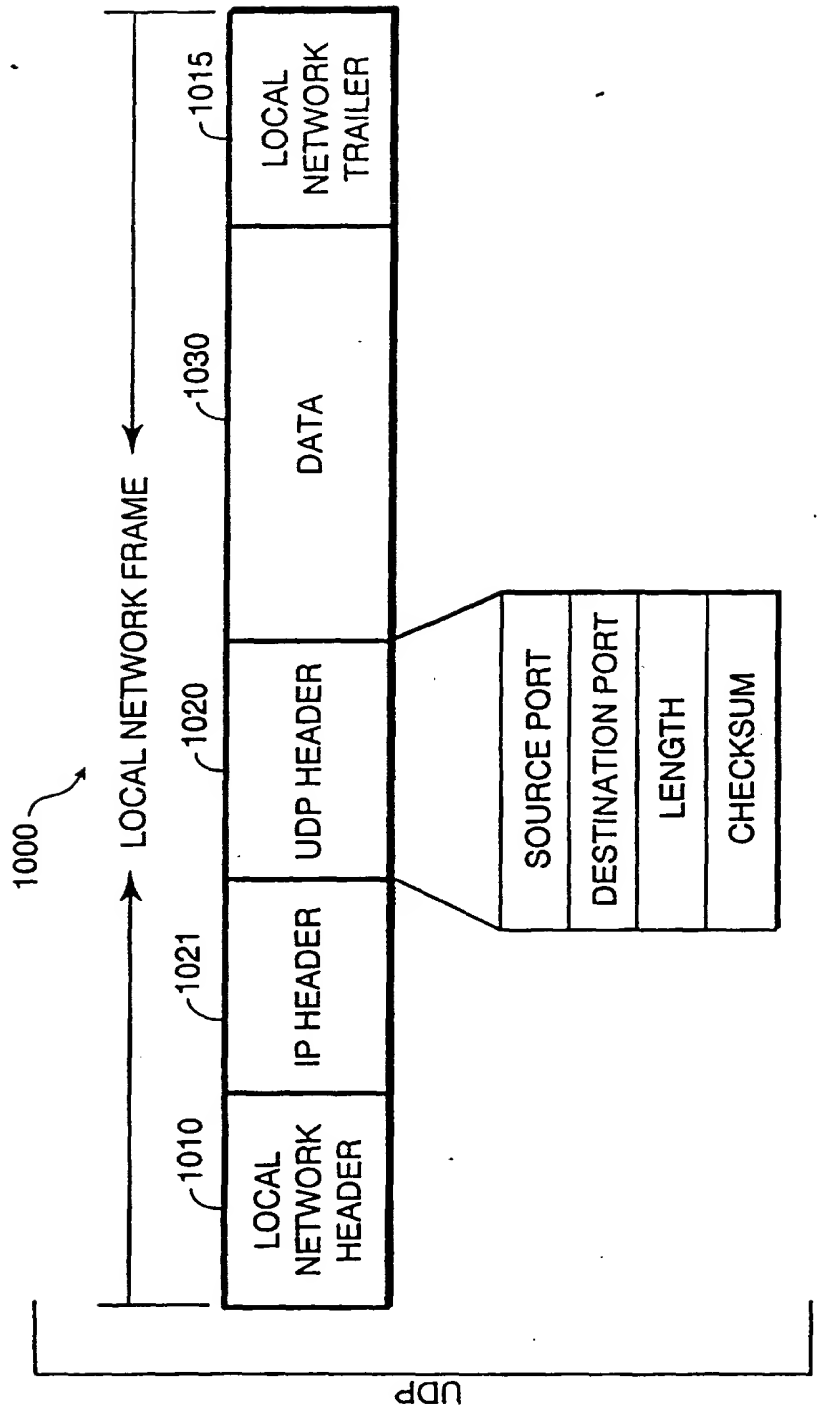


FIG. 10A

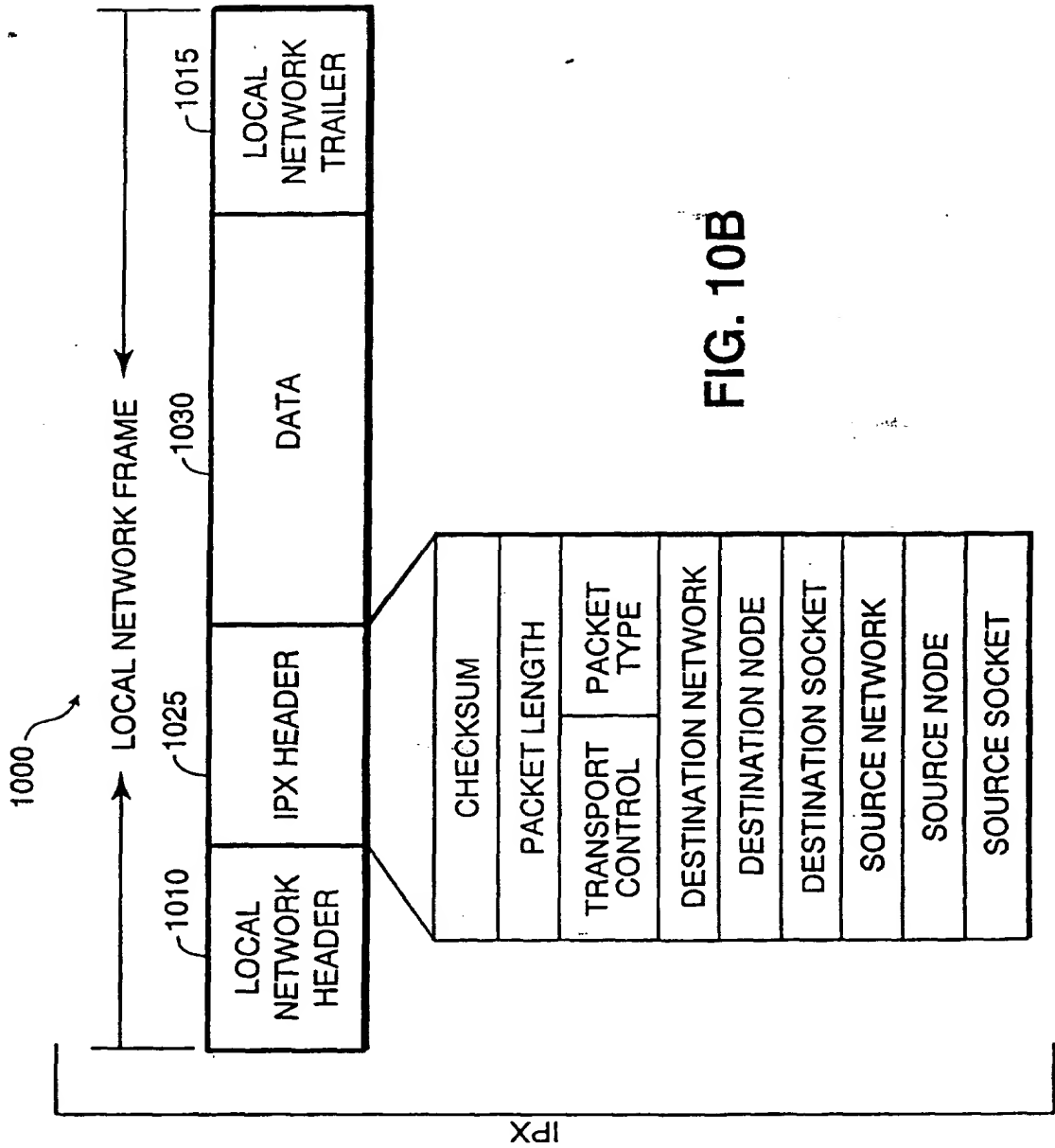


FIG. 10B

IPX

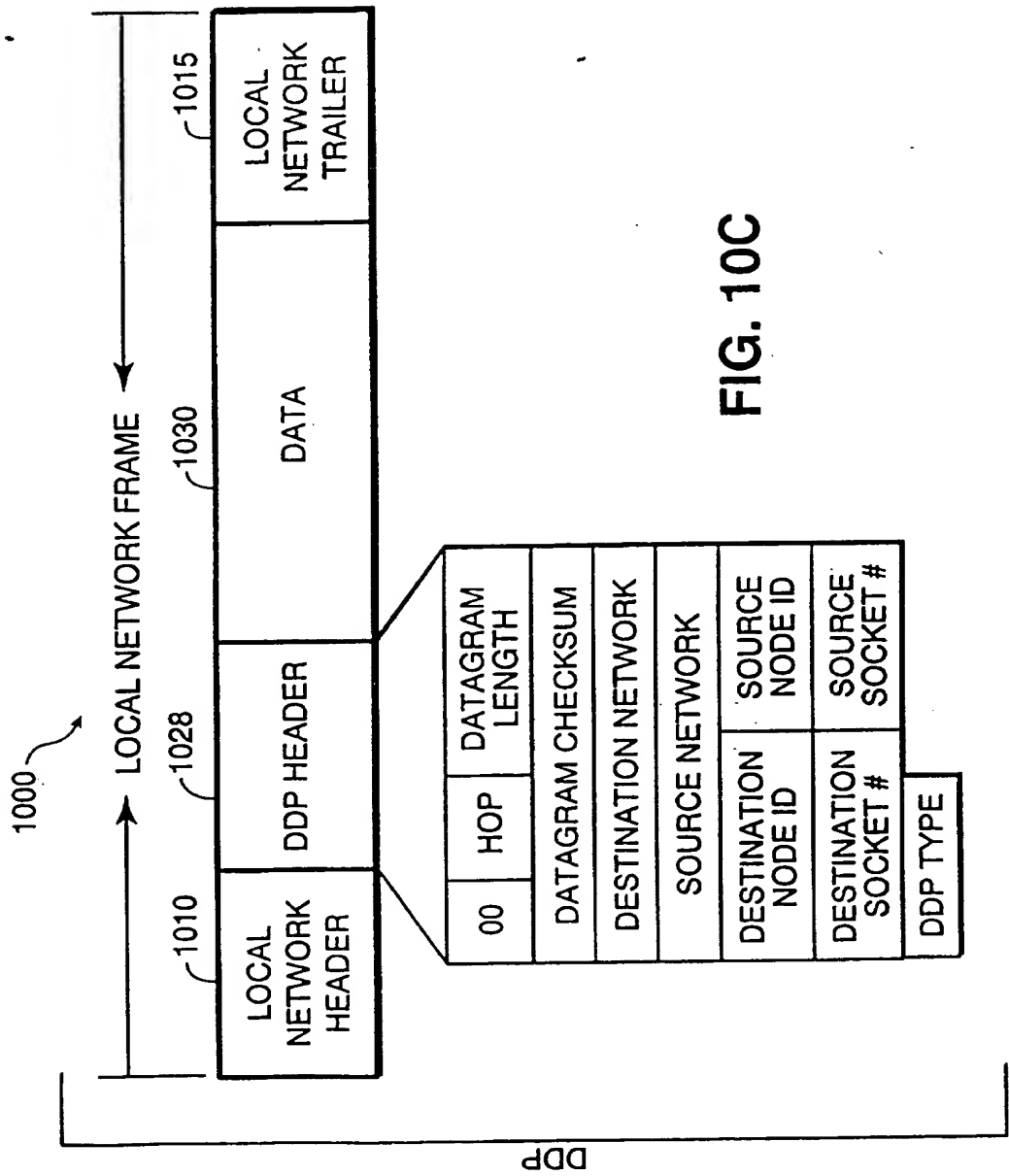


FIG. 10C

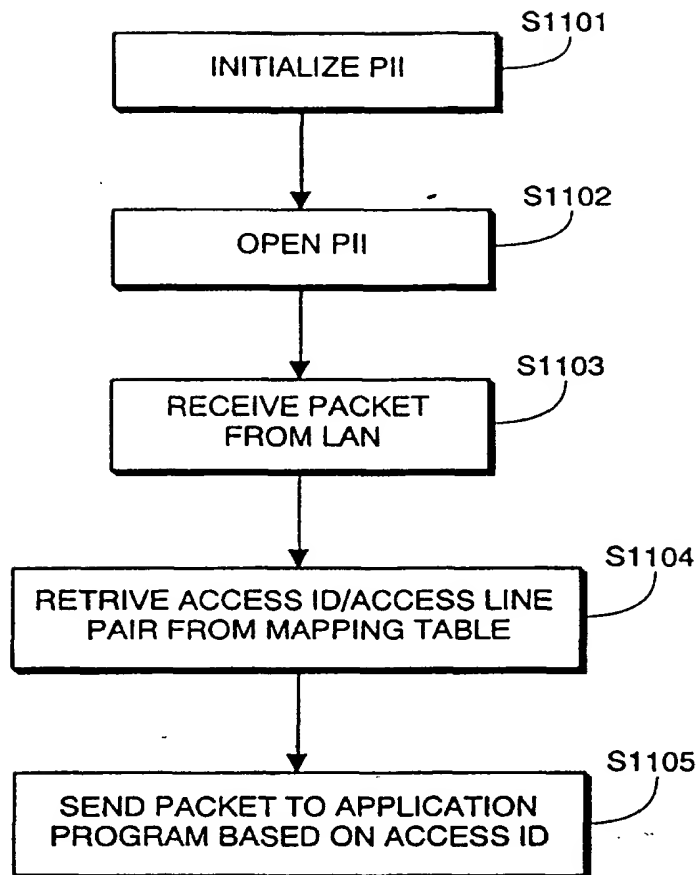


FIG. 11